

OpenMp

OMP Data Scoping

- Any variable that is invariant with respect to the parallel DO loop is a **scalar quantity** and if it is set in the parallel DO loop and then used it has to be private
- Any **scalar quantity** that is used and then set in the parallel DO Loop it gives rise to loop carried data recursion
- All Arrays referenced by the parallel loop index are shared
- Variables down the call chain revert to Fortran scoping rules
 - All COMMON, MODULE and SAVE statements have to be shared
 - All Automatic and allocated arrays have to be private
 - Data scoping down the call chain is EXTREMELY DIFFICULT

LESLIE3D OMP directives

C\$OMP PARALLEL DO

C\$OMP&SHARED (I1,CMU,GDIFFAC,DXI,RDENG,K2,DZI,KK,J2,CNUK,IBDD,XAREA)

C\$OMP&SHARED(K1,I2,DYI,J1,DTVOL,JMAX,DELTA,IADD,N,DCDX,DVDZ,CPK,FSI)

C\$OMP&SHARED(HFK,DWDY,QAV,WAV,CONC,PAV,EK,DU,VAV,HF,UAV,W,V,U,T,Q,H)

C\$OMP&PRIVATE (RHEV,RHEK,TXX,HK,ABD,CPAV,TEMP,SXX,SGSXX,WAVE,RLMBDA)

C\$OMP&PRIVATE (SGSEX,RHAVE,QSPI,VAVE,TZX,TXZ,RK,SGSXY,ICD,II,QSP,QX)

C\$OMP&PRIVATE (UAVE,SYY,SXZ,NS,RD,SGSXZ,TYX,TXY,EKAVE,IBD,SXY,DIV,SZZ)

C\$OMP&PRIVATE (L,J,I,QS,DVDY,DVDX,EKCOEF,DUDZ,DUDY,DUDX,DHDX,QDIFFX)

C\$OMP&PRIVATE (DKDX,RMU,DTDX,DWDZ,Dwdx,K)

DO K = K1,K2

DO J = J1,J2

ABD = DBLE(IBDD)

!

! EULER STUFF

!

DO I = I1,I2

QS(I) = UAV(I,J,K) * XAREA

END DO

IF(NSCHEME .EQ. 2) THEN

DO I = I1,I2

```
DO I = I1,I2
    DUDX(I) = DXI * (U(I+1,J,K) - U(I,J,K))
    DVDX(I) = DXI * (V(I+1,J,K) - V(I,J,K))
    DWDX(I) = DXI * (W(I+1,J,K) - W(I,J,K))
    DTDX(I) = DXI * (T(I+1,J,K) - T(I,J,K))
ENDDO
```

```
DO I = I1, I2
    DUDY(I) = DYI * (UAV(I,J+1,K) - UAV(I,J-1,K)) * 0.5D0
    DVDY(I) = DYI * (VAV(I,J+1,K) - VAV(I,J-1,K)) * 0.5D0
    DWDY(I) = DYI * (WAV(I,J+1,K) - WAV(I,J-1,K)) * 0.5D0
ENDDO
```

```
DO I = I1,I2
  FSI(I,2) = FSI(I,2) + SGSXX * XAREA
  FSI(I,3) = FSI(I,3) + SGSXY * XAREA
  FSI(I,4) = FSI(I,4) + SGSXZ * XAREA
  FSI(I,5) = FSI(I,5) + SGSEX * XAREA
END DO
```

```
DO I = I1,I2
  RDENG = EKCOEF(I) + RMU(I) / PRANDLT
  FSI(I,7) = FSI(I,7) - RDENG * DKDX(I) * XAREA
END DO
```

```
DO NS = 1,NSPECI
  DO I = I1,I2
    FSI(I,7+NS) = FSI(I,7+NS) -
>      EKCOEF(I) * DCDX(I,NS) * XAREA
  END DO
END DO
```

```

IF(ISGSK .EQ. 1) THEN
DO I = I1, I2
    RHAVE = 0.5D0 * (Q(I,J,K,1,N) + Q(I,J+1,K,1,N))
    EKAVE = 0.5D0 * (EK(I,J,K) + EK(I,J+1,K))
    EKCOEF(I) = RHAVE * CNUK * SQRT(EKAVE) * DELTA
    RHEV = 2.0D0 * EKCOEF(I)
!
    RDENG = CENAVE * EKCOEF(I)
    RHEK = TWO3 * RHAVE * EKAVE

    SXX = DUDX(I)
    SYY = DVDY(I)
    SZZ = DWDZ(I)
    SXY = 0.5D0 * (DUDY(I) + DVDX(I))
    SYZ = 0.5D0 * (DVDZ(I) + DWDY(I))

    DIV = (SXX + SYY + SZZ) * THIRD

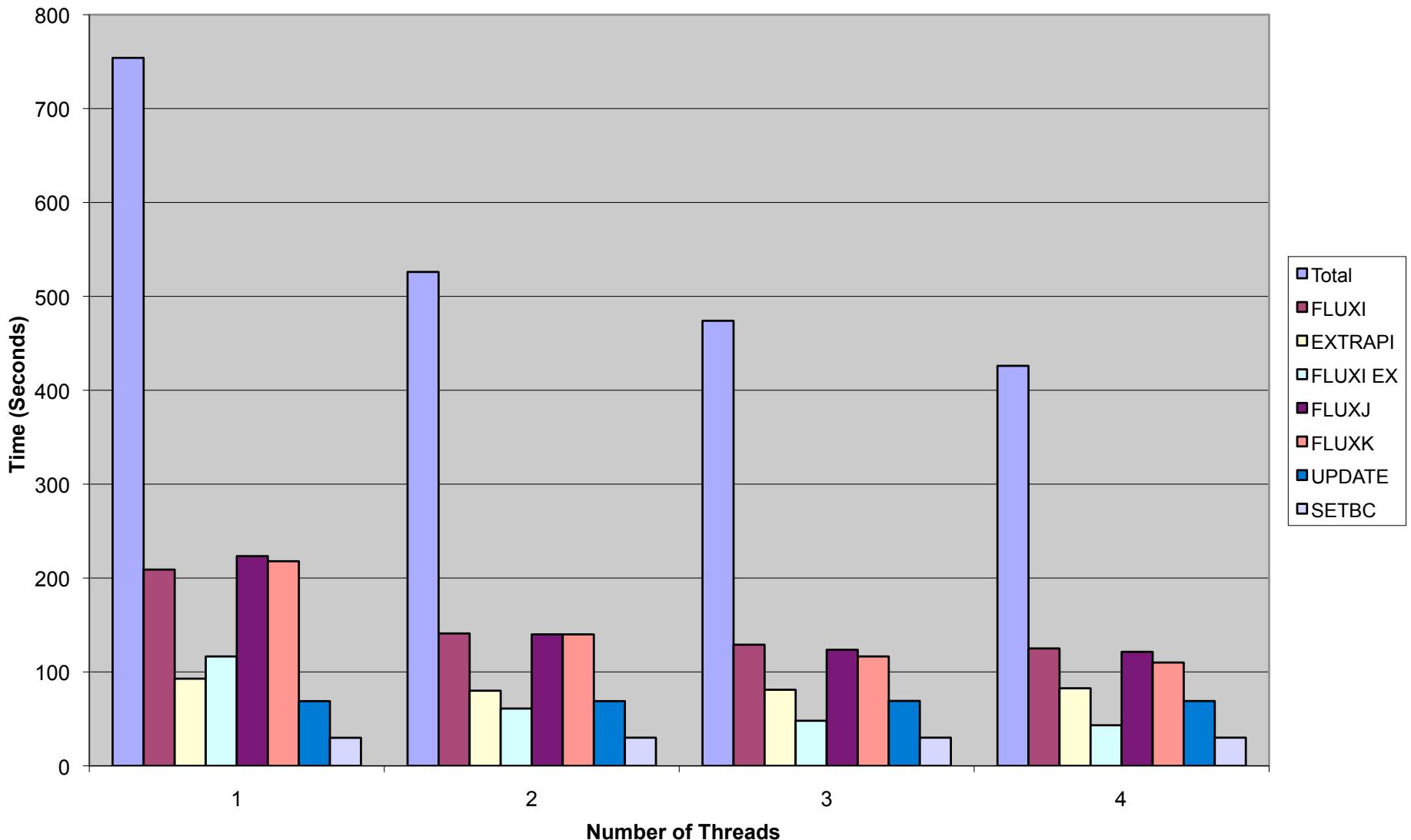
    SGSYY = - RHEV * (SYY - DIV) + RHEK
    SGSXY = - RHEV * SXY
    SGSYZ = - RHEV * SYZ

    SGSEY = - RDENG * DHDY(I)

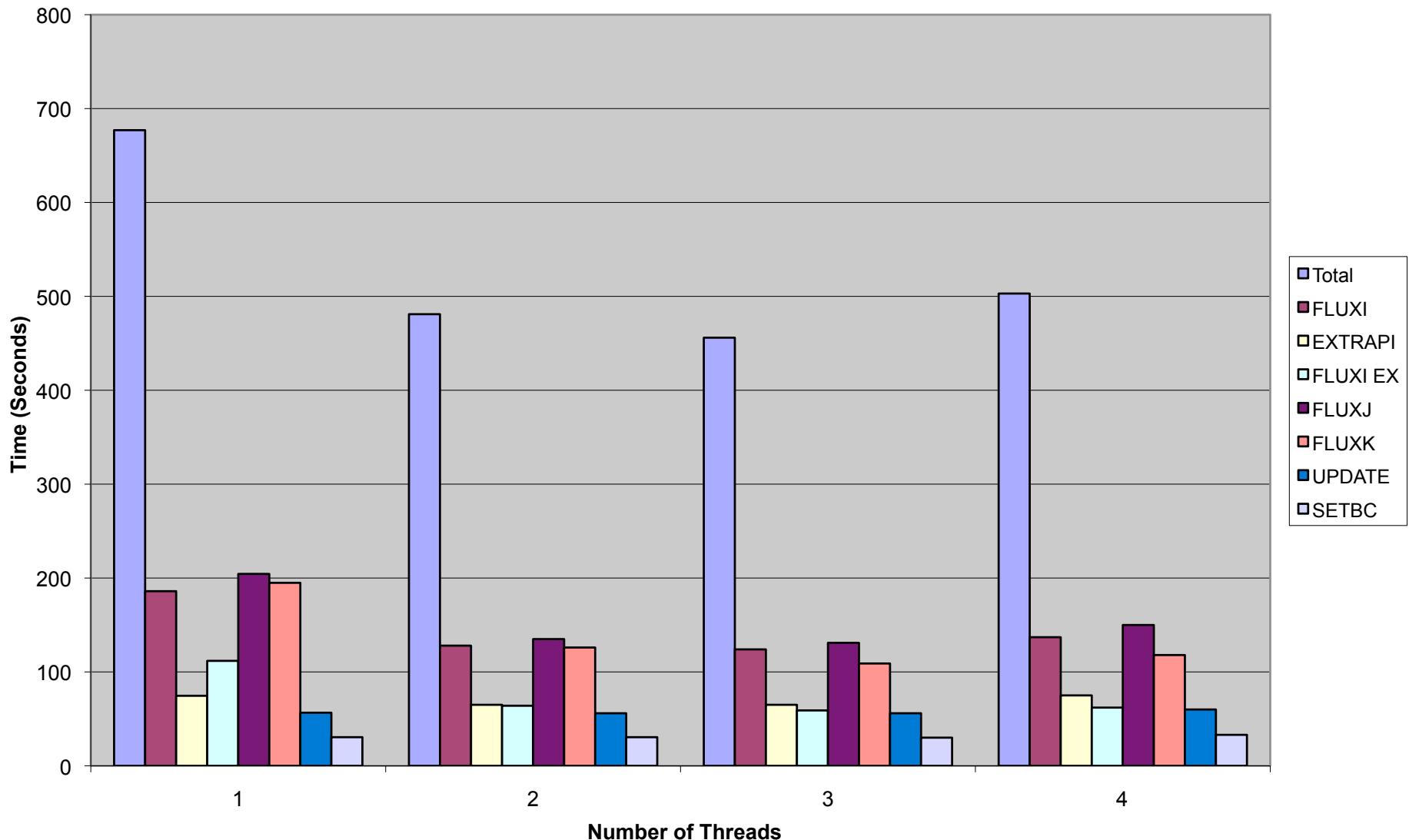
    FSJ(I,J,2) = FSJ(I,J,2) + SGSXY * YAREA
    FSJ(I,J,3) = FSJ(I,J,3) + SGSYY * YAREA
    FSJ(I,J,4) = FSJ(I,J,4) + SGSYZ * YAREA
    FSJ(I,J,5) = FSJ(I,J,5) + SGSEY
END DO

```

LESLIE 3D - Pathscale



leslie3d - PGI

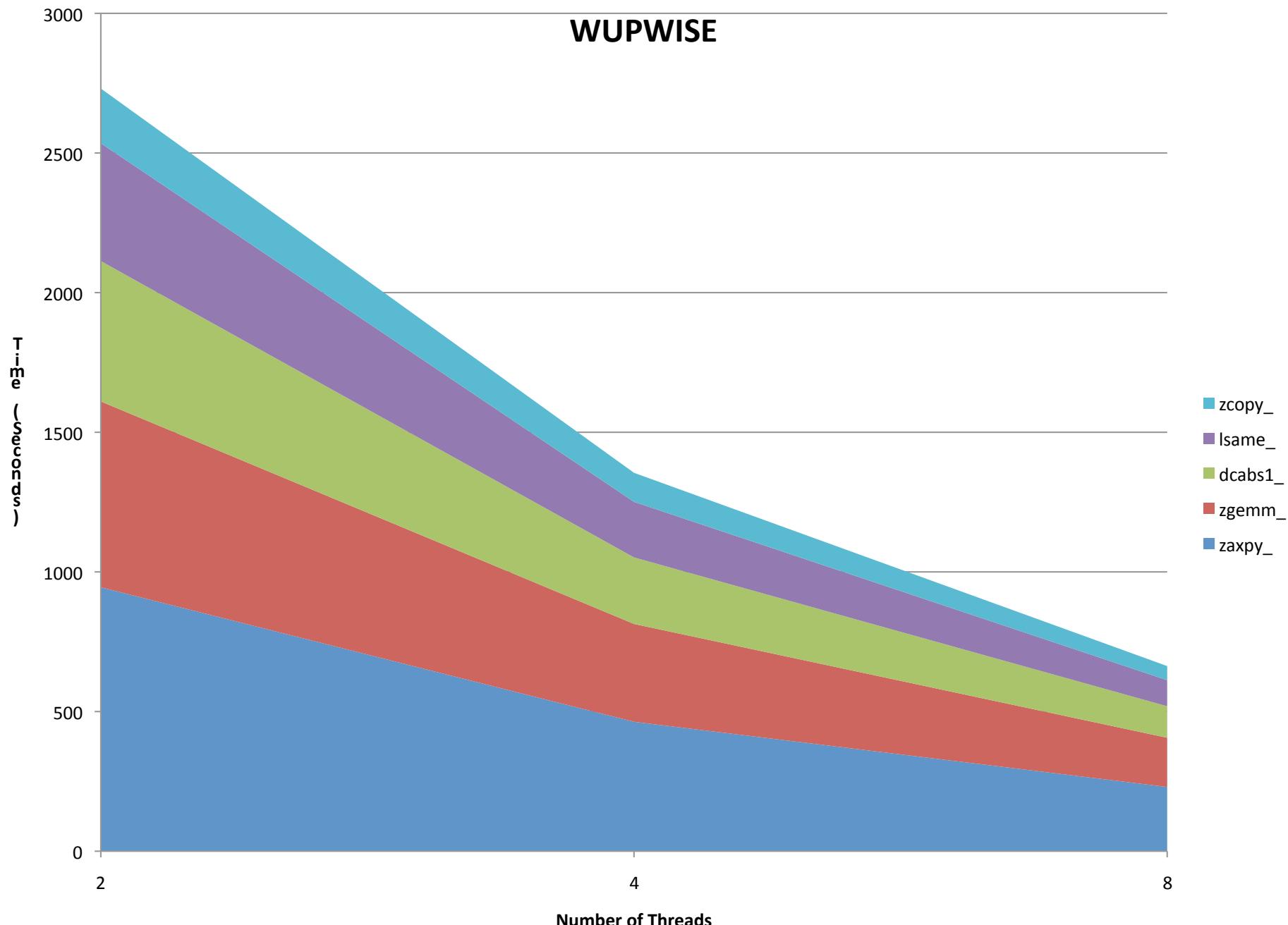


Performance Considerations

- Granularity of Computation
 - The parallel region should be as large, in number of operations, as possible
- Load Balancing
 - The work should be distributed evenly across the threads working the OpenMP region

SPEC OpenMP Routines

WUPWISE

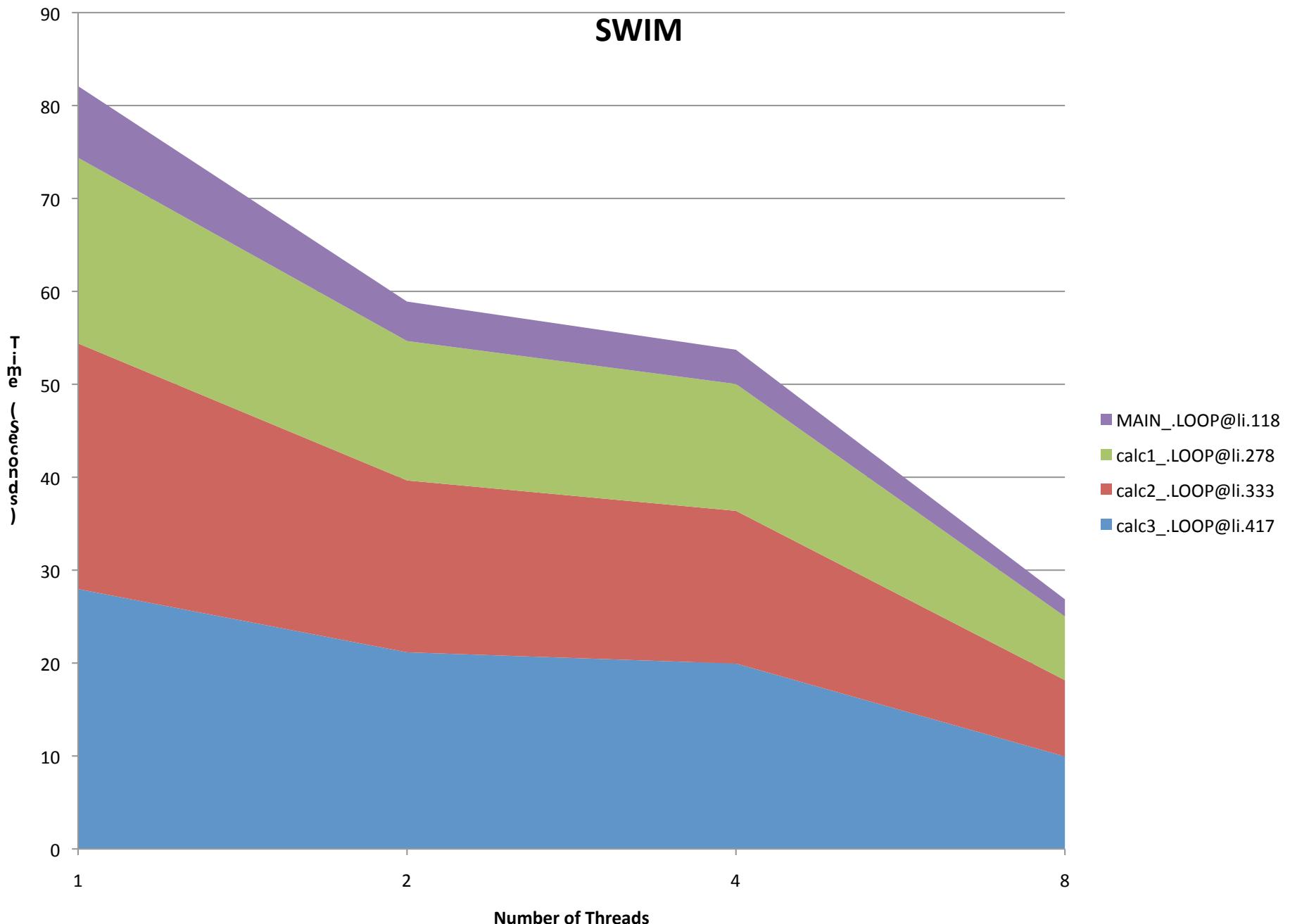


Major OMP Loop in WUPWISE

```
C$OMP PARALLEL
C$OMP+      PRIVATE (AUX1, AUX2, AUX3),
C$OMP+      PRIVATE (I, IM, IP, J, JM, JP, K, KM, KP, L, LM, LP),
C$OMP+      SHARED (N1, N2, N3, N4, RESULT, U, X)
C$OMP DO
    DO 100 JKL = 0, N2 * N3 * N4 - 1
        L = MOD (JKL / (N2 * N3), N4) + 1
        LP=MOD(L,N4)+1
        K = MOD (JKL / N2, N3) + 1
        KP=MOD(K,N3)+1
        J = MOD (JKL, N2) + 1
        JP=MOD(J,N2)+1
        DO 100 I=(MOD(J+K+L+1,2)+1),N1,2
            IP=MOD(I,N1)+1
            CALL GAMMUL(1,0,X(1,(IP+1)/2,J,K,L),AUX1)
            CALL SU3MUL(U(1,1,1,I,J,K,L),'N',AUX1,AUX3)
            CALL GAMMUL(2,0,X(1,(I+1)/2,JP,K,L),AUX1)
            CALL SU3MUL(U(1,1,2,I,J,K,L),'N',AUX1,AUX2)
            CALL ZAXPY(12,ONE,AUX2,1,AUX3,1)
            CALL GAMMUL(3,0,X(1,(I+1)/2,J,KP,L),AUX1)
            CALL SU3MUL(U(1,1,3,I,J,K,L),'N',AUX1,AUX2)
            CALL ZAXPY(12,ONE,AUX2,1,AUX3,1)
            CALL GAMMUL(4,0,X(1,(I+1)/2,J,K,LP),AUX1)
            CALL SU3MUL(U(1,1,4,I,J,K,L),'N',AUX1,AUX2)
            CALL ZAXPY(12,ONE,AUX2,1,AUX3,1)
            CALL ZCOPY(12,AUX3,1,RESULT(1,(I+1)/2,J,K,L),1)
100    CONTINUE
C$OMP END DO
```

Performance is Excellent

- Large Percentage of the code that uses the computation time is parallelized
- Granularity of the computation is very large
- Load Balance of the computation is good
- The computation that is parallelized is not memory bandwidth limited



Major OMP Loop in SWIM

```
!$OMP PARALLEL DO
DO 100 J=1,N
DO 100 I=1,M
CU(I+1,J) = .5D0*(P(I+1,J)+P(I,J))*U(I+1,J)
CV(I,J+1) = .5D0*(P(I,J+1)+P(I,J))*V(I,J+1)
Z(I+1,J+1) = (FSDX*(V(I+1,J+1)-V(I,J+1))-FSDY*(U(I+1,J+1)
1           -U(I+1,J)))/(P(I,J)+P(I+1,J)+P(I+1,J+1)+P(I,J+1))
H(I,J) = P(I,J)+.25D0*(U(I+1,J)*U(I+1,J)+U(I,J)*U(I,J)
1           +V(I,J+1)*V(I,J+1)+V(I,J)*V(I,J))
100 CONTINUE
```

Some Craypat Statistics

Time% 20.5%

Time 18.120602

Calls 120

PAPI_L1_DCM 232.754M/sec 4175600353 misses

PAPI_TLB_DM 1.334M/sec 23925543 misses

PAPI_L1_DCA 871.323M/sec 15631527182 refs

PAPI_FP_OPS 1932.840M/sec 34675154640 ops

User time (approx) 17.940 secs 41262000000 cycles 99.0%Time

Average Time per Call 0.151005 sec

CrayPat Overhead : Time 0.0%

HW FP Ops / User time 1932.840M/sec 34675154640 ops 21.0%peak(DP)

HW FP Ops / WCT 1913.576M/sec

Computational intensity 0.84 ops/cycle 2.22 ops/ref

MFLOPS (aggregate) 1932.84M/sec

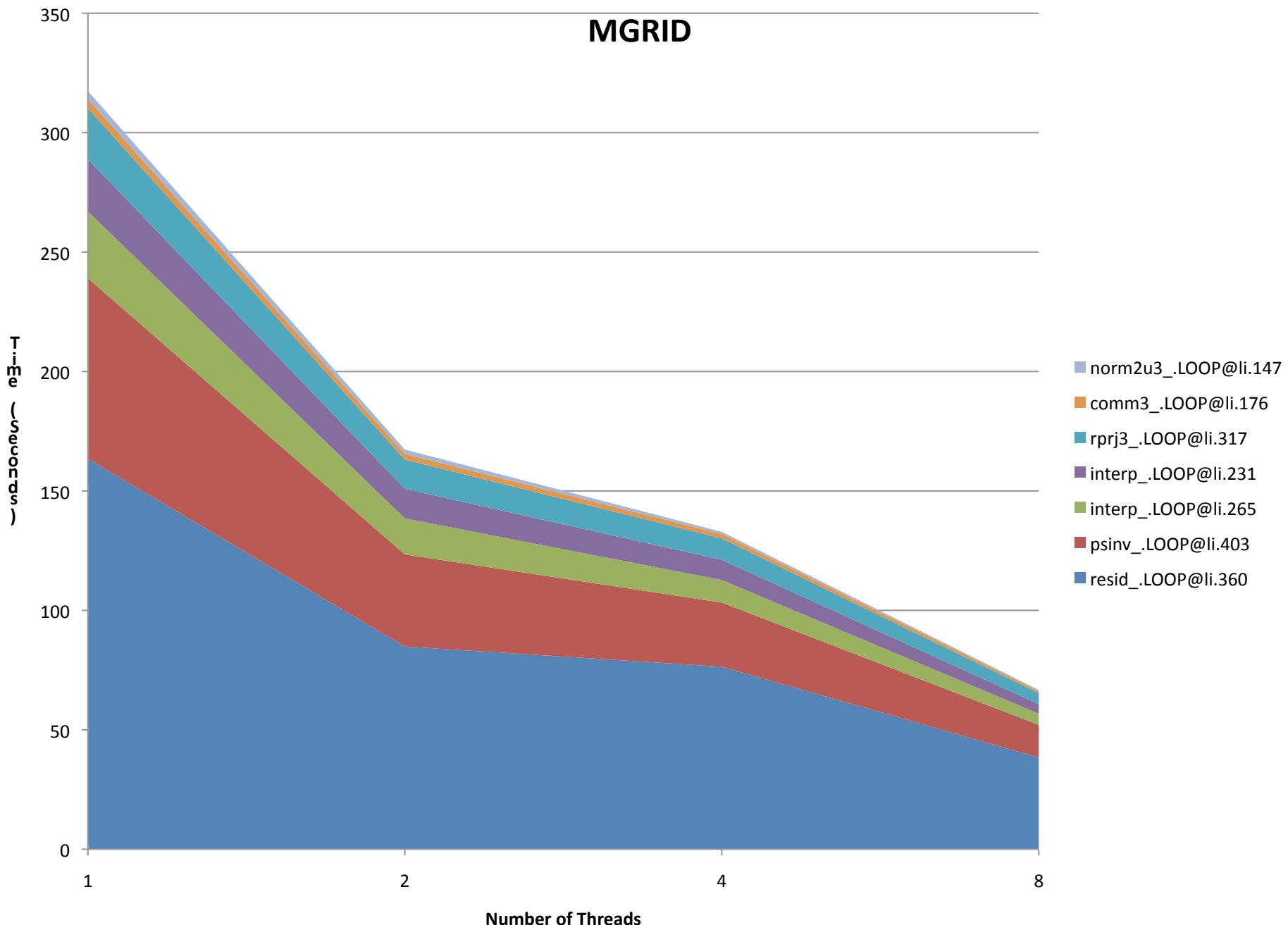
TLB utilization 653.34 refs/miss 1.276 avg uses

D1 cache hit,miss ratios 73.3% hits 26.7% misses

D1 cache utilization (M) 3.74 refs/miss 0.468 avg uses

Performance is Poor

- Large Percentage of the code that uses the computation time is parallelized
- Granularity of the computation is very large
- Load Balance of the computation is good
- The computation that is parallelized is memory bandwidth limited
 - Cache utilization is very bad so this is heavily dependent on memory loads/stores



Major OMP Loop in SWIM

MGRID is one of the NASA Parallel Benchmarks, it applies a multi-grid solver on a three dimensional grid. From the graph we once again see performance illustrative of a memory bound application.

In this case the Level 1 cache utilization is very good; however, the Level 2 cache utilization is very poor. Following are the derived metrics from hardware counters for Level 1 and 2 cache on the resid routine.

D1 cache hit,miss ratio 96.3% hits 3.7% misses

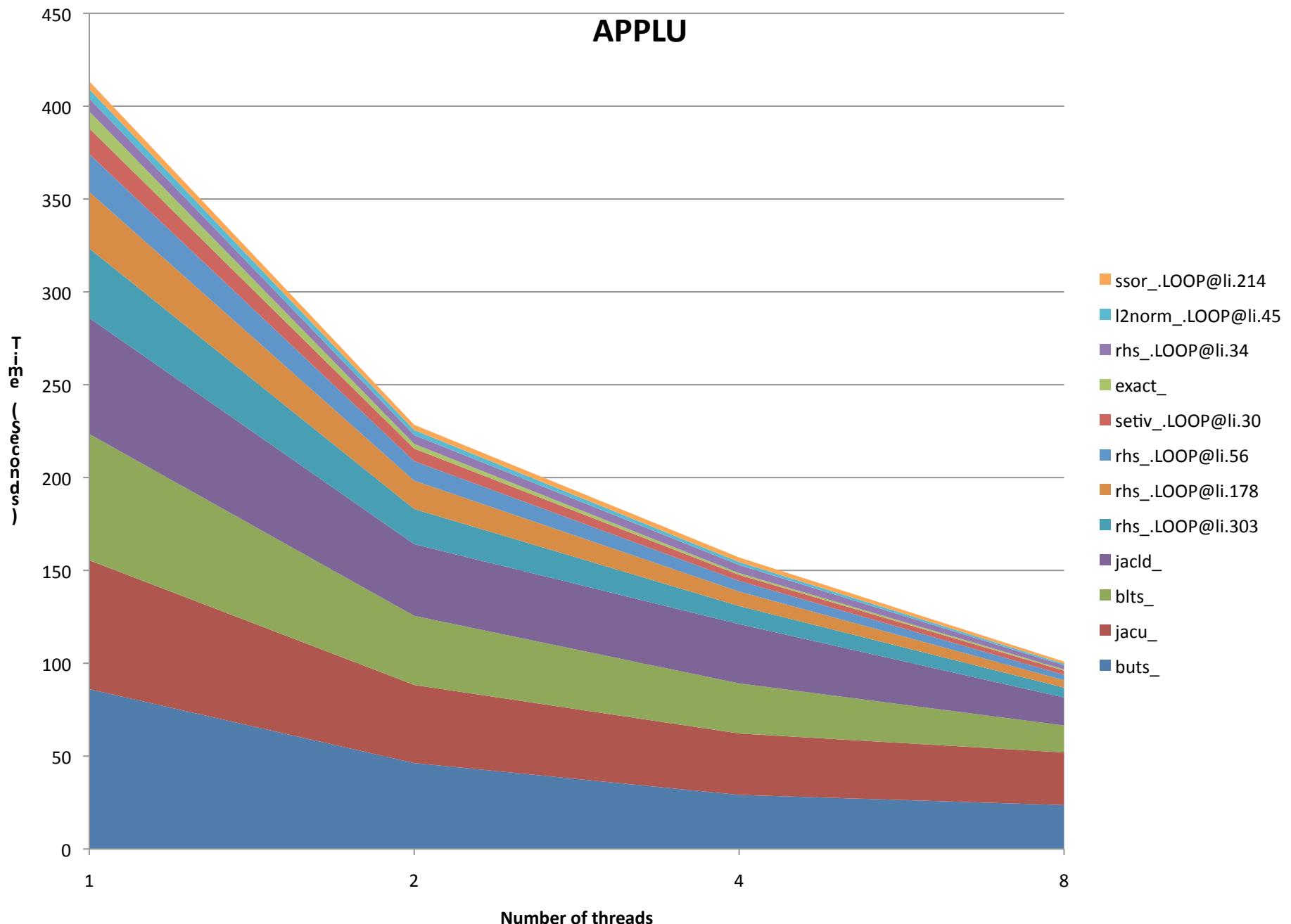
D2 cache hit,miss ratio 9.2% hits 90.8% misses

The important lesson is that poor cache utilization will steal OpenMP scalability, simply because it increases the reliance on memory bandwidth which is the rarest of all commodities on the node.

Performance is Poor

- Large Percentage of the code that uses the computation time is parallelized
- Granularity of the computation is very large
- Load Balance of the computation is good
- The computation that is parallelized **is very** memory bandwidth limited
 - Cache utilization is very bad so this is heavily dependent on memory loads/stores

APPLU

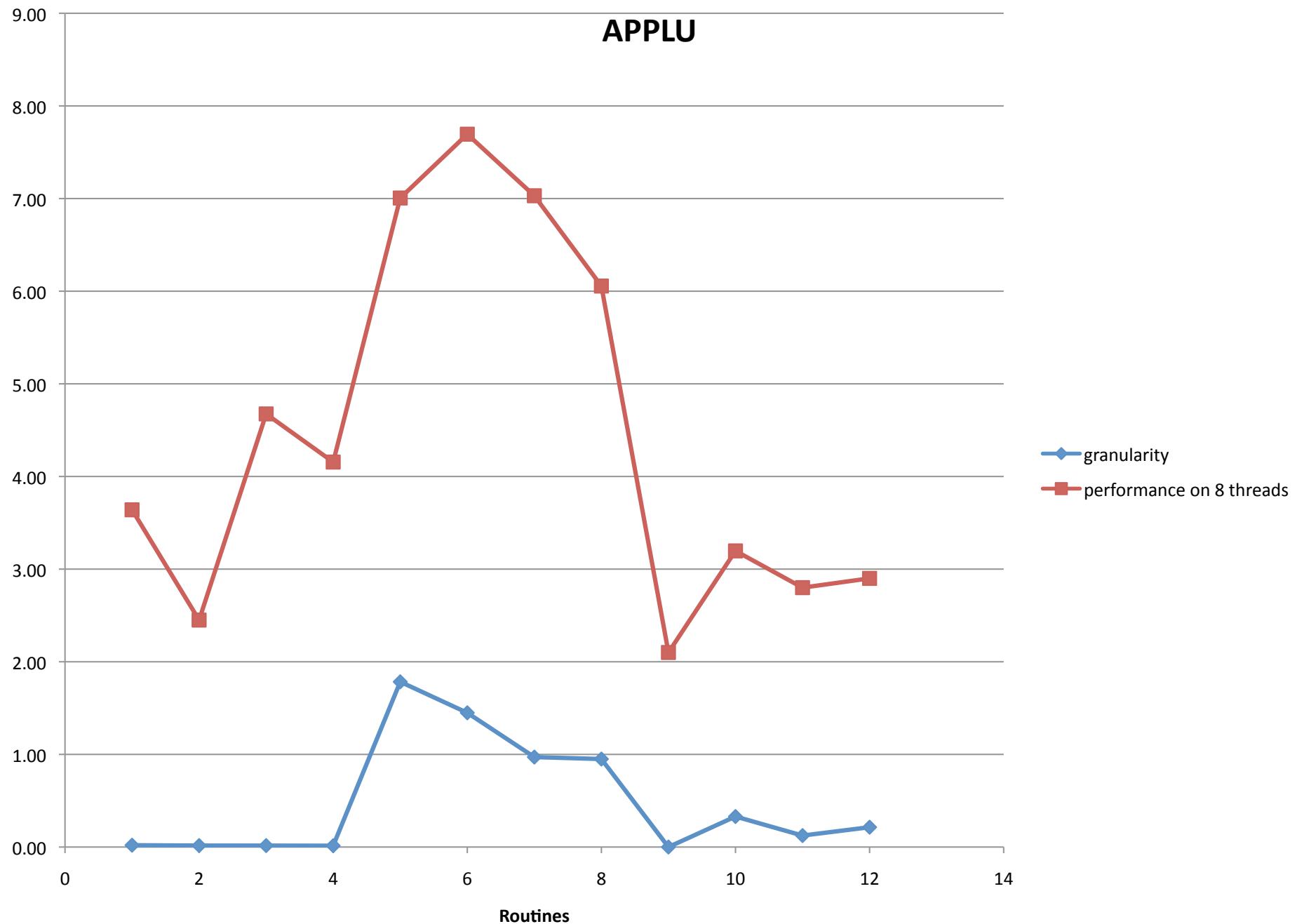


APPLU is another NASA parallel benchmark which performs the solution of five coupled nonlinear PDE's, on a 3-dimensional logically structured grid, using an implicit psuedo-time marching scheme, based on two-factor approximate factorization of the sparse Jacobian matrix. The following chart shows that some routines scale very well while other do not. The overall performance is 3.91 on 8 threads. The reason for the difference in the performance of the individual routines can be attributed to the granularity of the parallelized region. The table below shows the granularity and performance gain for each of 7 of the major routines.

	buts	jacu	blts	jacld	Rhs_303	Rhs_153	Rhs_56
Granularity	.02	.02	.02	.01	1.78	1.45	.97
Performance	3.6	2.45	4.67	4.15	7	1.69	7.03

While it is not a linear function of granularity, the performance is highly dependent upon the granularity of the loop. In this case granularity is the time the loop takes to execute divided by the number of times the loop is executed. The other variation in the table could be due to the memory bandwidth requirement of the individual loops.

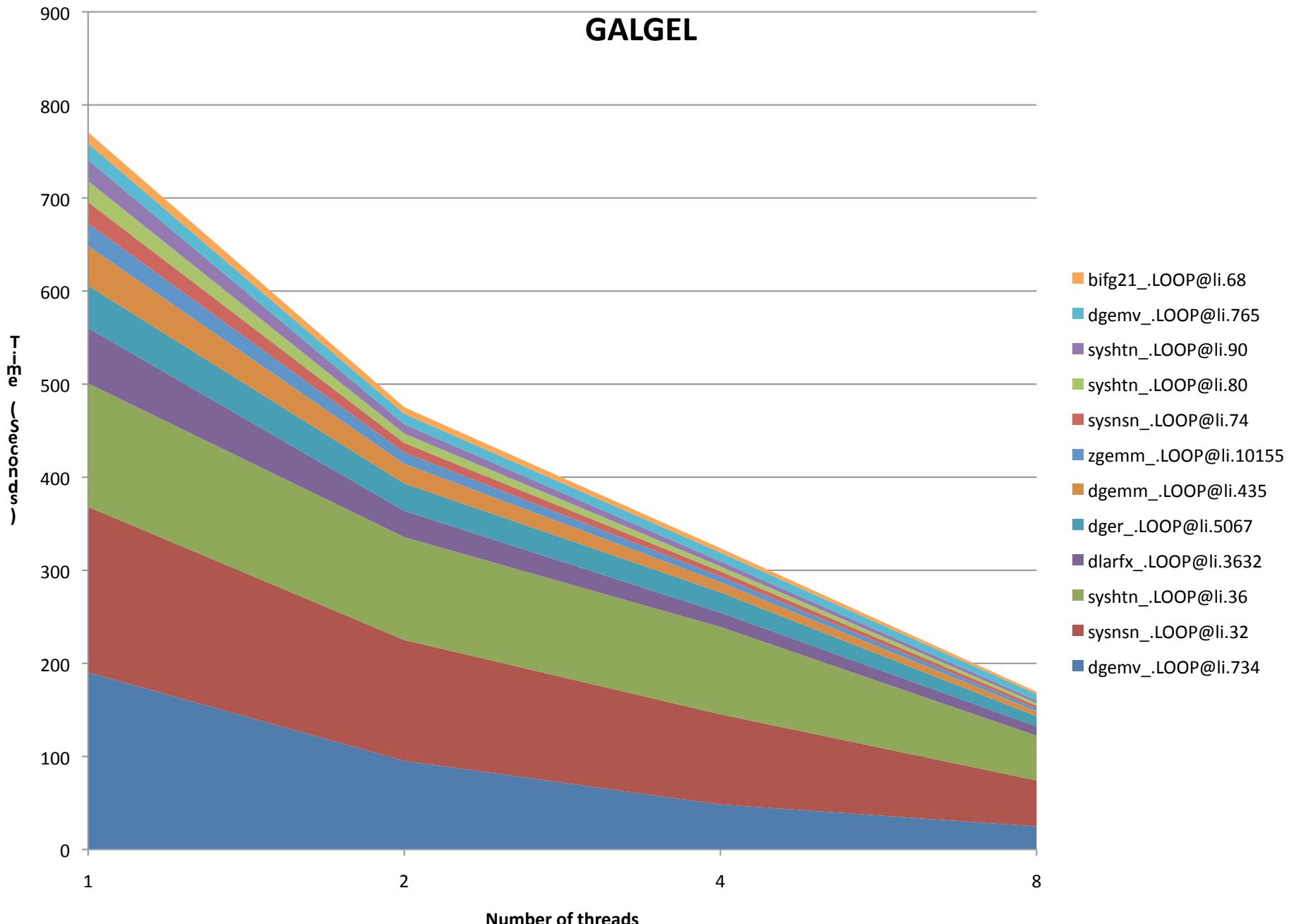
APLU



Performance is So-so

- Large Percentage of the code that uses the computation time is parallelized
- Granularity of the computation is **not good**
- Load Balance of the computation is good
- The computation that is parallelized **is** memory bandwidth limited

GALGEL

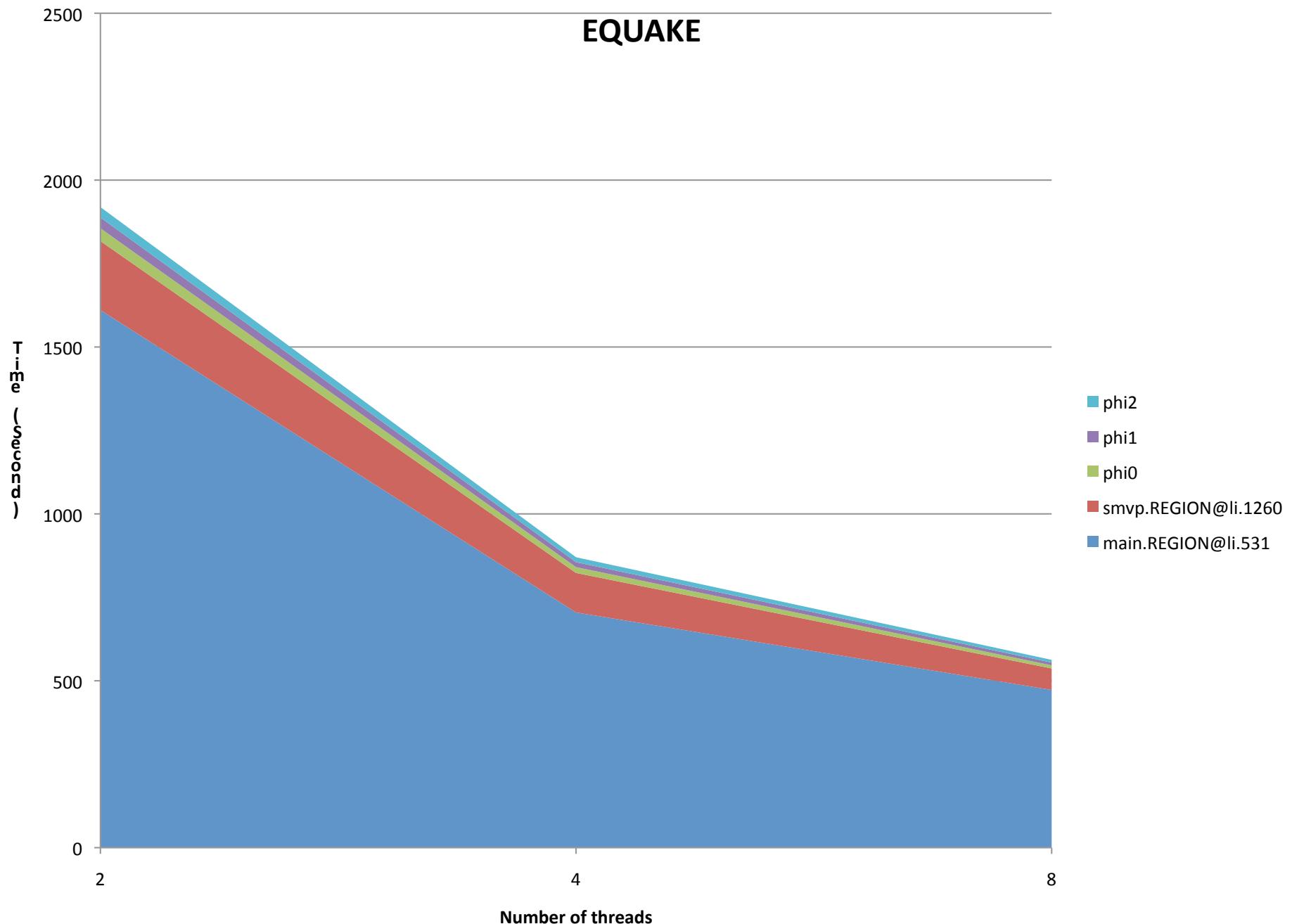


Major Loop in GALGEL

```
!$OMP DO SCHEDULE(GUIDED)
    Ext12: Do LM = 1, K
        L = (LM - 1) / NKY + 1
        M = LM - (L - 1) * NKY
        Do IL=1,NX
            Do JL=1,NY
                Do i=1,NKX
                    Do j=1,NKY
                        LPOP( NKY*(i-1)+j, NY*(IL-1)+JL ) = &
                            WXTX(IL,i,L) * WXTY(JL,j,M) + &
                            WYTX(IL,i,L) * WYTY(JL,j,M)
                    End Do
                End Do
            End Do
        End Do
        LPOP1(1:K) = MATMUL( LPOP(1:K,1:N), Y(K+1:K+N) )
        Proj3( NKY*(L-1)+M, 1:K) = LPOP1(1:K)
        Xp(NKY*(L-1)+M) = DOT_PRODUCT (Y(1:K), LPOP1(1:K) )
        Proj4( NKY*(L-1)+M,1:N) = &
            MATMUL( TRANSPOSE( LPOP(1:K,1:N) ), Y(1:K) )
    End Do Ext12
 !$OMP END DO
```

Performance is So-so

- Large Percentage of the code that uses the computation time is parallelized
- Granularity of the computation is **not good**
- Load Balance of the computation is good
- The computation that is parallelized **is very** memory bandwidth limited



```

#pragma omp for
for (i = 0; i < nodes; i++) {
    Anext = Aindex[i];
    Alast = Aindex[i + 1];

    sum0 = A[Anext][0]*v[i][0] + A[Anext][0][1]*v[i][1] + A[Anext][0][2]*v[i][2];
    sum1 = A[Anext][1][0]*v[i][0] + A[Anext][1][1]*v[i][1] + A[Anext][1][2]*v[i][2];
    sum2 = A[Anext][2][0]*v[i][0] + A[Anext][2][1]*v[i][1] + A[Anext][2][2]*v[i][2];

    Anext++;
    while (Anext < Alast) {
        col = Acol[Anext];

        sum0 += A[Anext][0]*v[col][0] + A[Anext][0][1]*v[col][1] +
            A[Anext][0][2]*v[col][2];
        sum1 += A[Anext][1][0]*v[col][0] + A[Anext][1][1]*v[col][1] +
            A[Anext][1][2]*v[col][2];
        sum2 += A[Anext][2][0]*v[col][0] + A[Anext][2][1]*v[col][1] +
            A[Anext][2][2]*v[col][2];

        if (w2[my_cpu_id][col] == 0) {
            w2[my_cpu_id][col] = 1;
            w1[my_cpu_id][col].first = 0.0;
            w1[my_cpu_id][col].second = 0.0;
            w1[my_cpu_id][col].third = 0.0;
        }

        w1[my_cpu_id][col].first += A[Anext][0][0]*v[i][0] + A[Anext][1][0]*v[i][1] +
            A[Anext][2][0]*v[i][2];
        w1[my_cpu_id][col].second += A[Anext][0][1]*v[i][0] + A[Anext][1][1]*v[i][1] +
            A[Anext][2][1]*v[i][2];
        w1[my_cpu_id][col].third += A[Anext][0][2]*v[i][0] + A[Anext][1][2]*v[i][1] +
            A[Anext][2][2]*v[i][2];
        Anext++;
    }

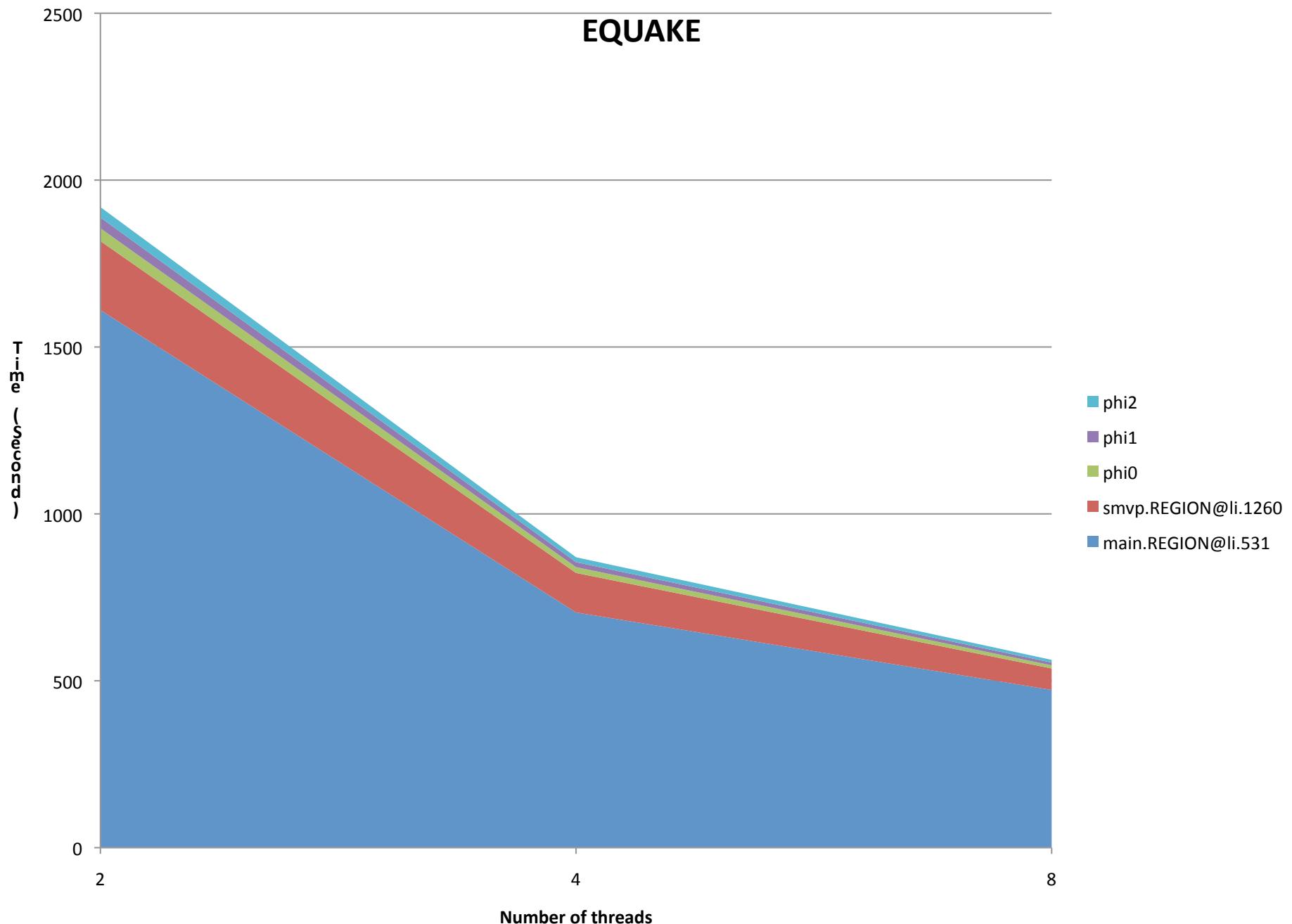
    if (w2[my_cpu_id][i] == 0) {
        w2[my_cpu_id][i] = 1;
        w1[my_cpu_id][i].first = 0.0;
        w1[my_cpu_id][i].second = 0.0;
        w1[my_cpu_id][i].third = 0.0;
    }

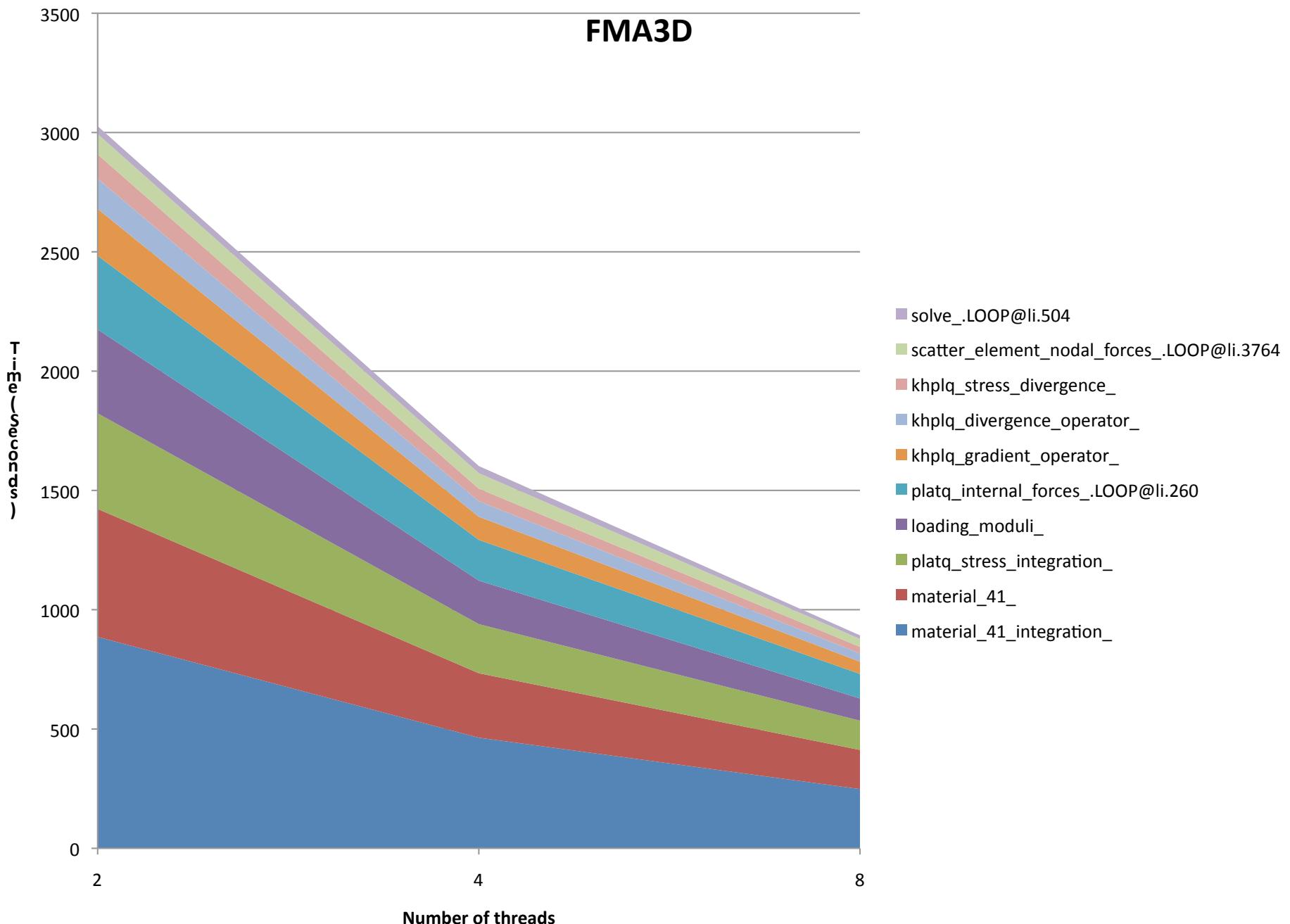
    w1[my_cpu_id][i].first += sum0;
    w1[my_cpu_id][i].second += sum1;
    w1[my_cpu_id][i].third += sum2;
}
}

```

Performance is Good

- Large Percentage of the code that uses the computation time is parallelized
- Granularity of the computation is **questionable**
- Load Balance of the computation is good
- The computation that is parallelized is not memory bandwidth limited





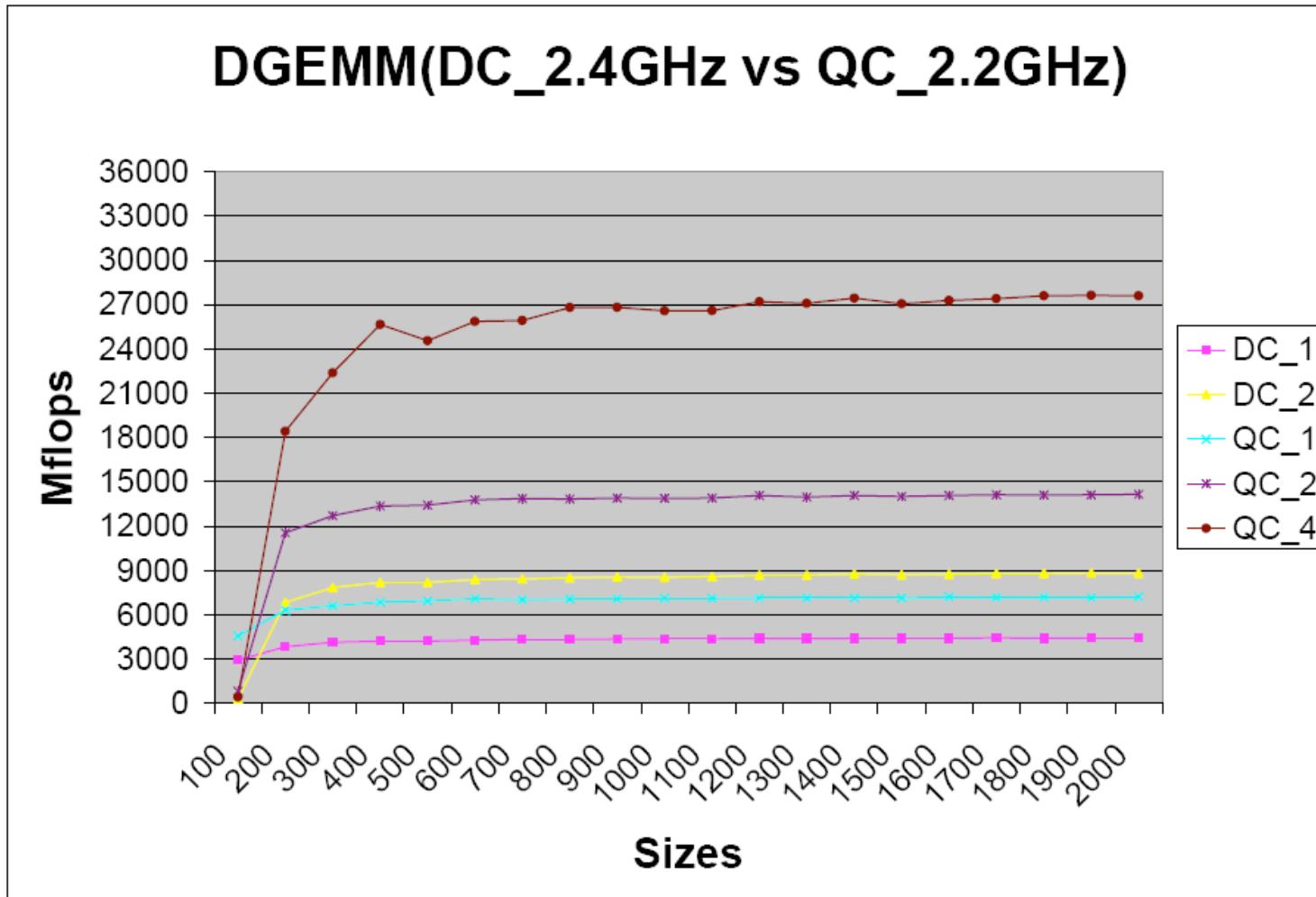
```
!$OMP PARALLEL DO DEFAULT(SHARED) PRIVATE(N,IX)
DO N = 1,NUMP4
  IF (PLATQ(N)%PAR%IGR .EQ. CURRENT_GROUP) THEN
    IX(1:8) = PLATQ(N)%PAR%IX(1:8)
    FORCE(IX(1:8))%Xint = FORCE(IX(1:8))%Xint + PLATQ(N)%RES%Xint
    FORCE(IX(1:8))%Yint = FORCE(IX(1:8))%Yint + PLATQ(N)%RES%Yint
    FORCE(IX(1:8))%Zint = FORCE(IX(1:8))%Zint + PLATQ(N)%RES%Zint
  ENDIF
ENDDO
!$OMP END PARALLEL DO
```

Indirect address causing poor memory utilization – only achieve 2.5 in going from 2 to 8 threads

Performance is Good

- Large Percentage of the code that uses the computation time is parallelized
- Granularity of the computation is good
- Load Balance of the computation is good
- The computation that is parallelized **is** memory bandwidth limited

OpenMP in LIBSCI



OpenMP in LIBSCI

